On the Minimal Problems of Low-Rank Matrix Factorization

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Abstract

Low-rank matrix factorization is an essential problem in many areas including computer vision, with applications in e.g. affine structure-from-motion, photometric stereo, and non-rigid structure from motion. However, very little attention has been drawn to minimal cases for this problem or to using the minimal configuration of observations to find the solution. Minimal problems are useful when either outliers are present or the observation matrix is sparse. In this paper, we first give some theoretical insights on how to generate all the minimal problems of a given size using Laman graph theory. We then propose a new parametrization and a building-block scheme to solve these minimal problems by extending the solution from a small sized minimal problem. We test our solvers on synthetic data as well as real data with outliers or a large portion of missing data and show that our method can handle the cases when other iterative methods, based on convex relaxation, fail.

1. Introduction

Assume that we have a matrix $X \in \mathbb{R}^{m \times n}$ of rank $r < \min(m, n)$. Low-rank matrix factorization aims at finding two matrices $U \in \mathbb{R}^{r \times m}$ and $V \in \mathbb{R}^{r \times n}$, such that $X$ can be written as

$$ X = U^T V. \quad (1) $$

In geometric computer vision, several problems can be formulated as low-rank matrix factorization. For example, for affine structure-from-motion (SfM), the observation matrix containing feature tracks can be factorized into the camera motions and the 3D structure. For photometric stereo, the directions of light sources and the surface normals are separated by factorizing the measurement matrix composed of pixel intensities under a Lambertian model. More examples of applications can be found in [2, 3, 11, 30].

For real data, with noise or outliers, the observation matrix will in general not have the exact rank as given. So the problem is usually formulated as minimizing the following cost function,

$$ \min_{U,V} \|X - U^T V\|, \quad (2) $$

where $\|\cdot\|$ is a matrix norm – typically the $L_1$- or $L_2$-norm. This is under the assumption that all entries in $X$ are known. When missing data is present, a binary matrix $W \in \{0, 1\}^{m \times n}$ is used to indicate if a certain entry $X(i,j)$ is present ($W(i,j) = 1$) or missing ($W(i,j) = 0$). In this case, the problem is formulated as

$$ \min_{U,V} \|W \odot (X - U^T V)\|, \quad (3) $$

where $\odot$ is the Hadamard product, that is, the element-wise product.

Related work  Early work on low-rank matrix factorization solve (2) using the $L_2$-norm. Truncating the singular value decomposition of $X$ has been shown to give the optimal solution under the $L_2$-norm when $X$ is complete, see [9]. The work in [29] was the first to consider missing data.

However, the $L_2$-norm is sensitive to outliers. An early work that aims for robustness to outliers is [1]. Here they use an iteratively re-weighted least squares approach to optimize the objective function. In [4], a damped Newton method is proposed to handle the missing data. Since [16], more robust norms have been considered. In [16], algorithms based on alternating optimization are introduced under the Huber-norm and the $L_1$-norm. In [10], the Wiberg algorithm [29] is generalized from the $L_2$-norm to the $L_1$-norm. A number of recent works put extra constraints on the factor matrices $U$ and $V$. In [5], the constraints that the matrix $U$ should lie on a certain manifold for different applications are considered and incorporated in the formulation. In [31], orthogonal constraints on the columns of $U$ and a nuclear norm regularizer on $V$ are incorporated.

Most methods mentioned above are based on alternating optimization and are prone to get trapped in local minima. Recently, several works [8, 7, 21, 11, 25] re-formulate the problem to minimize the convex surrogate of the rank
function, that is, the nuclear norm. This makes it possible to use convex optimization to find the global optimum of the approximated objective function. These approaches can handle the problems when the rank is not known a priori. However, for applications with a given rank, the nuclear norm based methods usually perform inferior to the bilinear formulation-based methods [6]. Another drawback of the nuclear norm method is that it is sensitive to outliers. The performance is also affected by an increasing number of missing data.

A few recent works [23, 15, 20, 19] also explore the idea to divide the whole matrix into overlapping sub-blocks and combine the sub-block solutions. One motivation is due to the structured missing data patterns in the measurement matrix. For example, for SfM datasets the observations normally span a diagonal band while off-diagonal entries are mainly missing data. However, these methods do not consider both outliers and missing data at the same time. We use a similar block-partition strategy, but give a more general solution by considering both outliers and missing data.

In this paper, we investigate the minimal problems in low-rank matrix factorization. It turns out that there are many of them and that the structure of these minimal problems is extremely rich and complicated. Nevertheless, we will give some results concerning characterization and generation of such problems and also provide an initial study on solving them. It is worth noting that a few works are relevant to this paper. For the specific SfM problem, estimation of the minimal cases for missing data were investigated in [26], but here we look at a much more general problem. In [14], a combinatorial method is proposed by searching among the linearly solvable minimal cases, while in this paper, we provide a unified view on all the minimal problems and give solvers for more general minimal problems, not only restricted to linearly solvable ones. The concepts based on rigid graphs have also been previously investigated for multidimensional scaling problems with missing data in [27].

Our main contributions in this paper are: (i) generation and characterization of the minimal problems, (ii) solvers to a number of these minimal problems and (iii) algorithms using minimal solvers to factorization problem with missing data and outliers in $L_1$-norm or $L_2$ norm.

2. Generating the minimal problems

We can define the following partial order relationship "≤" between two index matrices $W$ and $W'$.

**Definition 2.1.** Given two index matrices $W$ and $W'$, we say that $W \leq W'$ if $W(i,j) = 1 \implies W'(i,j) = 1$. Here $W$ is submatrix of $W'$. $W < W'$ if $W$ is a strict submatrix of $W'$.

An index matrix $W$ is said to be rigid if for general data, the low-rank matrix factorization problem given by

$$W \odot (X - U^TV) = 0,$$

is locally well defined.

The notion of rigidity is invariant under the permutation of the rows or columns. We define the following equivalent relationship between two index matrices

**Definition 2.2.** Given two index matrices $W$ and $W'$, we say that $W$ is equivalent to $W'$ if $W(i,j) = W'(P_1(i), P_2(j))$ where $P_1$ and $P_2$ are permutations of row/column indices for all $\{i, j\}$.

Before giving the definition of the minimal problems in matrix factorization, we look at the degrees of freedom (DoF) for a low-rank matrix. For a rank-$r$ factorization of a matrix $X \in \mathbb{R}^{m \times n}$ we have that $U$ has $mr$ DoF and $V$ has $nr$ DoF. There is a total coordinate ambiguity of size $r \times r$ as

$$X = U^TV = U^TQQ^{-1}V,$$

where $Q \in \mathbb{R}^{r \times r}$ is a full rank matrix. Thus a matrix $X \in \mathbb{R}^{m \times n}$ with rank $r$, has $mr + nr - r^2$ degrees of freedom. This means that we need at least $d = mr + nr - r^2$ measurements to recover a rank-$r$ matrix $X$ of size $m \times n$.

A minimal problem for low-rank matrix factorization is characterized using a minimal index matrix, which is defined as

**Definition 2.3.** An index matrix $W$ for a rank-$r$ problem is said to be minimal if it is rigid and satisfies $\sum_{i,j} W(i,j) = mr + nr - r^2$. $W$ is said to be overdetermined if $\sum_{i,j} W(i,j) > mr + nr - r^2$.

It is trivial to see that if $W$ is rigid and if $W \leq W'$ then $W'$ is also rigid. It also can be shown that if $W'$ is rigid and overdetermined, then there is at least one $W < W'$ that is rigid and minimal.

A minimal low-rank matrix factorization problem is finding two factor matrices $U$ and $V$ that exactly solves (4), where $W$ is a minimal index matrix and $X$ is the measurement matrix. For the minimal problem, with general coefficients, characterized by a minimal index matrix $W$, there is a finite number $n_W > 0$ of solutions, where $n_W$ only depends on the index matrix $W$.

2.1. Characterizing minimal index matrices

For every minimal index matrix $W$ of size $m \times n$ there is a corresponding minimal index matrix $W'$ of size $n \times m$ such that $W(i,j) = W'(j,i)$. Without loss of generality we may thus in the discussion assume that $n \geq m$. Each minimal index matrix has $mr + nr - r^2$ non-zero elements. Since by assumption $m \leq n$, we have at most $2mr - r^2$ non-zero elements, which are distributed among $n$ columns. Thus there are never enough non-zero elements to fill up $2r$-
non-zero elements of each column, in another word, there is at least one column which has fewer than $2r$ non-zeros. Furthermore it is obvious that for rank-$r$ problems, the minimal index matrices must have at least $r$ non-zero elements in each column, otherwise the corresponding column $v$ of $V$ has too few constraints to be solvable. So for the column with the smallest number $k$ of non-zero elements we must have $r \leq k < 2r$.

We also notice that since we assume $r < \min(m,n)$ then $m-r \geq 1$. This means that for a minimal index matrix, the number of non-zeros $nr + nr - r^2 \geq nr + r$. If we distribute these non-zeros in $n$ columns, each of which has at least $r$ zeros, then there is at least one column which has $k$ non-zeros, where $k \geq r + 1$.

2.2. Henneberg-like extensions

We will now describe how to generate the minimal problems in low-rank matrix factorization. The inspiration comes from rigid graph theory, where the Henneberg construction is used to generate the Laman graph, see [21, 12]. The idea is that one starts with the smallest index matrix, and by a series of extensions every index matrix can be generated. For example, for $r = 2$, the smallest index matrix is

$$W = \begin{bmatrix} 1 & 1 \\ 1 & 1 \end{bmatrix}. \quad (6)$$

In the following we will distinguish between constructive extensions and non-constructive extensions. For a constructive extension from $W$ to $W'$, one can infer the number of solutions $n_{W'}$ from $n_W$ and construct the solver, denoted by $f_{W'}$ from $f_W$. For non-constructive extensions, it can be shown that $W$ is minimal if and only if $W'$ is minimal. However, we can in general neither infer the number of solutions $n_{W'}$ from $n_W$ nor derive a solver $f_{W'}$ from $f_W$. We propose the following extensions and reductions which are denoted Henneberg-$k$ extensions/reductions. Of these Henneberg-1 is constructive, whereas Henneberg-$k$ are generally non-constructive.

Henneberg-1 extension Given a minimal index matrix $W$ for a rank-$r$ problem of size $m \times n$, an extended minimal index matrix $W'$ of size $m \times (n+1)$ is formed by adding a column with exactly $r$ indices (non-zero elements). The number of solutions is identical, i.e. $n_{W'} = n_W$. Extending an algorithm from $f_W$ to $f_{W'}$ is straightforward. A similar extension can be done by adding a row with $r$ indices.

Henneberg-1 reduction Given a minimal index matrix $W$ for the rank-$r$ problem of size $m \times n$ where there is a column $w_j$ with exactly $r$ indices, a reduced minimal index matrix $W'$ of size $m \times (n-1)$ is formed by removing column $j$. The number of solutions is preserved under the reduction and the solution for $W'$ can be obtained straightforward from the solution for $W$. A similar extension can be done by removing a row with $r$ indices.

Henneberg-2 extension Given a minimal index matrix $W$ for the rank-$r$ problem of size $m \times n$, where there is a column $w_j$ with at least $r+1$ non-zero elements at rows $i_1, \ldots, i_{r+1}$ (such a column must exist from Section 2.1), an extended index matrix $W'$ of size $m \times (n+1)$ can be formed by first adding a column $w'$ with exactly $r+1$ non-zero elements at rows $i_1, \ldots, i_{r+1}$, then setting one of the non-zeros of $w_j$ to be zero. The resulting index matrix $W'$ will be minimal. A similar extension can be done in a row-wise manner.

Similarly one can define the Henneberg-2 Reduction. This can generalized to Henneberg-$k$ extension and reduction for $k > 2$. We illustrate an example of how the index matrices for rank-2 problem can be generated using both Henneberg-1 and Henneberg-2 extensions in Fig. 1.

As shown in [12], every minimally rigid graph can be formed using a sequence of Henneberg-1 and Henneberg-2 construction in the context of rigid graphs. In the following we will show that every minimal index matrix can be generated using a series of Henneberg-$k$ extension defined above for rank $r = 1$ and $r = 2$ case. We also make a conjecture that this is the case for the general rank-$r$ problems with $r > 2$.

Theorem 2.4. Each index matrix for a minimal rank-1 problem can be formed by a series of Henneberg-1 extensions from the $1 \times 1$ index matrix $W = [1]$ as the base case.

Proof. The proof is by induction over size. If the matrix is of size $1 \times 1$ then we are finished. Assume that it is true for all matrices of size $m \times n$ with $m + n \leq K$. Take a minimal index matrix with $m + n = K + 1$. From Section 2.1, we know that there always exists a column with at least $k$ non-zero elements where $r \leq k < 2r$. In this case, there is always a column with exactly one non-zero element. After a Henneberg-1 reduction on that column, we obtain a minimal index matrix with $m + n = K$, which can be constructed as our assumption. So the original matrix is a Henneberg-1 extension of a smaller index matrix that is in the assumption, which proves the theorem. A similar proof exists for a row-wise extension. \qed
**Theorem 2.5.** Each index matrix for a minimal rank-2 problem can be formed by a series of Henneberg-1 and Henneberg-2 extensions from the $2 \times 2$ index matrix in (6) as the base case.

**Proof.** The proof is similarly by induction over size. If the matrix is of size $2 \times 2$ then we are finished. Assume that it is true for all matrices of size $m \times n$ with $m + n \leq K$. Take a minimal index matrix with $m + n = K + 1$. From Section 2.1, we know that there is always a column with exactly $k$ non-zero elements where $r \leq k < 2r$, i.e. $k = 2$ or $3$ in this case. If the column has two non-zeros then the index matrix can be constructed using the Henneberg-1 extension from the one that is in the assumption. If the column has three non-zeros then it can be reduced to an index matrix of size $m + n \leq K$ using a Henneberg-2 reduction. In either case, we have shown that the index matrix can be constructed using a Henneberg-1 or Henneberg-2 extension from a smaller index matrix given in the assumption, which proves the theorem. The proof for a row-wise extension can be shown similarly. 

**Conjecture 2.1.** Each minimal index matrix for the rank $r$ problem can be formed by a series of Henneberg-1 to Henneberg-$r$ extensions.

Thus it is possible to generate all the minimal index matrices using a sequence of Henneberg extensions.

**3. Minimal solvers**

In this section, we will describe the solvers to the minimal problems generated from different Henneberg extensions.

**3.1. Solvers for Henneberg-1 extensions**

Suppose we are given a rank-$r$ minimal problem with index matrix $W$ and measurements $X$. Assume that the solution to $\{X, W\}$ is given by $\{U, V\}$ as

$$W \odot (X - U^TV) = 0. \quad (7)$$

Now we apply a Henneberg-1 extension to $W$ to get $W' = [W | w]$ where the column vector $w$ has $r$ ones. Correspondingly the measurement matrix is extended to $X' = [X | x]$

where $x$ has $r$ observations. To find the solution to the extended minimal problem, it is obvious that we only need to solve for the extra column $v$ of $V$ to satisfy the following equation

$$W' \odot (X' - U^T [V | v]) = 0. \quad (8)$$

Now assume the positions for the $r$ observations of $x$ are at $I = \{n_1, \ldots, n_r\}$. Then we have the following equation

$$U^T_I v = x_I, \quad (9)$$

where $U_I \in \mathbb{R}^{r \times r}$ and $x_I \in \mathbb{R}^r$ denotes taking the corresponding rows at $I$ from $U$ and $x$ respectively. This is a linear system with a unique solution. For the row-wise Henneberg-1 extension, we will keep $V$ unchanged and solve for the extra row $u^T$ of $U$ using a similar strategy.

**3.2. Solvers for other constructive extensions**

Henneberg-$k$ extensions for $k \geq 2$ are non-constructive, which means that for $W \rightarrow W'$ one cannot construct the solution to $W'$ from the solution to $W$. However, we can define some other constructive extensions. The intuitive idea is that given two index matrices $W_1$ and $W_2$, one can construct a new index matrix $W$ by ”glueing” $W_1$ and $W_2$ together. By ”glueing”, we mean that $W$ contains both $W_1$ and $W_2$ with overlapping rows and columns as illustrated in Fig. 3.

In the following, we will first describe a general parametrization for these types of constructive extensions, that is independent of the rank $r$. We derive several constraints from the parametrization that can be used to solve the minimal problem. A few examples for rank-2 and rank-3 are then illustrated.

**Parametrization** Consider that we have a minimal problem with index matrix $W$ and measurement $X$. $W$ is constructed by ”glueing” two index matrices $W_1$ and $W_2$ as in Fig. 3a. We use $I_1$ and $J_1$ to denote the row and column indices of $W_1$ in $W$, similarly $I_2$, $J_2$ for $W_2$. Then we have
$W_1 = W(I_1, J_1)$ and $W_2 = W(I_2, J_2)$. Accordingly for measurement matrices $X_1$ and $X_2$ associated with $W_1$ and $W_2$, we have $X_1 = X(I_1, J_1)$ and $X_2 = X(I_2, J_2)$. We also use $I_{12}$ and $J_{12}$ to denote the indices of overlapping rows and columns as $I_{12} = I_1 \cap I_2$ and $J_{12} = J_1 \cap J_2$.

Assume that the solutions to the sub-problem $\{W_1, X_1\}$ and $\{W_2, X_2\}$ are given by $\{U_1, V_1\}$ and $\{U_2, V_2\}$ respectively. To construct the solution to $\{W, X\}$, the idea is to find a transformation matrix $H \in \mathbb{R}^{r \times r}$ to transform the subspace $U_2$ to the same coordinate framework as the subspace $U_1$. Using this transformation we have

$$U_2^TV_2 = (U_2)^TH^TH^{-T}(V_2) = (HU_2)^T(H^{-T}V_2). \quad (10)$$

Now $HU_2$ and $H^{-T}V_2$ are in the same coordinate framework as $U_1$ and $V_1$ respectively. The remaining problem is to solve for $H$. We have the following constraint, that states that $U_1$ and $HU_2$ should coincide for the overlapping columns as

$$U_1(i, I_{12}) = HU_2(i, I_{12}), \quad i = 1, \ldots, r. \quad (11)$$

Similarly we have the overlapping constraints for $V_1$ and $H^{-T}V_2$ as

$$V_1(i, J_{12}) = H^{-T}V_2(i, J_{12}), \quad i = 1, \ldots, r. \quad (12)$$

This can be written as

$$H^TV_1(i, J_{12}) = V_2(i, J_{12}), \quad i = 1, \ldots, r, \quad (13)$$

which is linear with respect to $H$.

If we have enough constraints from (11) and (13), $H$ can be solved linearly. In that case, the overlap should be of size $r \times r$ as in Fig. 3a. For the cases where the overlap doesn’t give sufficiently many constraints, we need some extra constraint outside $W_1$ and $W_2$ to solve for the transformation matrix $H$ as in Fig. 3b-c.

To solve the case with extra constraints, we know that each extra measurement $X(i, j)$ directly gives a constraint as

$$u_i^TV_j - X(i, j) = 0, \quad (14)$$

where $u_i$ and $v_j$ are $i$-th and $j$-th column of $U$ and $V$ respectively. We need to express the constraint using $U_1$, $U_2$, $V_1$ and $V_2$ from which we assume $\{U_1, V_1\}$ are in the same coordinate system as $\{U, V\}$. There are generally two cases. When $X(i, j)$ is in the top-right as in Fig. 3b, one needs to use $U_1$ and $V_2$ as

$$U_1(:, i_1)^TH^{-T}V_2(:, J_2) - X(i, j) = 0. \quad (15)$$

where $i_1$ denotes the local index of column $i$ in $U_1$ and $j_2$ the local index of column $j$ in $V_2$. When $X(i, j)$ is in the bottom-left as in Fig. 3c, one needs to rewrite (14) using $U_2$ and $V_1$ instead as

$$(HU_2(:, i_2))^TV_1(:, J_1) - X(i, j) = 0. \quad (16)$$

where $i_2$ denotes the local index of column $i$ in $U_2$ and $j_1$ the local index of column $j$ in $V_1$.

All these constraints from (11), (13), (15) and (16) form either a linear system or a simple polynomial system, from which one can solve the transformation $H$ and thus get a solution to the original problem. In the following we will illustrate a few examples for rank-2 and rank-3 problem, all of which can be parametrized and solved in a similar way.

Examples of rank-2 constructive extensions Here we will present a few constructive extensions for rank-2 problems, which are illustrated in Fig. 4. The overlap for rank-2 problems is at most $2 \times 2$, otherwise the extension is non-minimal. For a $2 \times 2$ overlap, $H$ can be solved linearly using only the overlap constraints. For a $1 \times 1$ overlap with an extra constraint as in Fig. 4a, the overlap in $U$ gives two equations from (11) and the overlap in $V$ gives two equations from (13). Among the four equations, one is redundant as it is automatically satisfied when the other three hold. So one can parametrize $H$ using a single variable, and we can solve using a single extra constraint.

One could of course generate a minimal problem by “glueing” more index matrices. In Fig. 4b-c, each sub-part needs a transformation, namely $H_1$, $H_2$ and $H_3$ of which one can fix the gauge by setting $H_2 = I$. Two extra constraint are needed to solve the problem. For the case in Fig. 4d, four transformation $H_1$, $H_2$, $H_3$ and $H_4$ are needed. By setting $H_2 = I$, one can parametrize $H_1$, $H_3$ and $H_3^{-1}H_4$ using $z_1$, $z_2$ and $z_3$ respectively. Each of three extra constraints provide a quadratic equation in $z_1$, $z_2$ and $z_3$ respectively, which could be solved giving three solutions.

Examples of rank-3 constructive extensions For rank-3 problems, the maximum overlap is $3 \times 3$ which can then be solved linearly. The $2 \times 2$ overlap case, illustrated in Fig. 5a, $H$ can be parametrized using a single variable $z$, and can be solved linearly using one extra constraint. When the overlap is $1 \times 1$ as in Fig. 5b, the overlap provides three equations in (11) and three equations in (13). However one of them is redundant, which means it reduces the DoF of $H$ from nine to four. Thus four extra constraints are needed to solve for $H$.

4. Algorithms for structured data patterns

For real problems, a measurement matrix might be more dense than that for a minimal problem. However, one could apply a random sampling strategy similar to [14] to sample minimal configurations from measurements and solve a minimal problem in each RANSAC-like iteration.

In applications where the locations of the missing data are highly correlated and structured – for example
5. Experiments

We have conducted a number of experiments on both synthetic and real data. For synthetic data and affine structure-from-motion, we use the solvers for both the Henneberg-1 extension and the constructive extensions from Section 3.2. The other constructive extension is useful when the measurement matrix is sparse, especially when it is not possible to find a minimal configuration using only the Henneberg-1 extension. For the shape basis estimation, we use only the Henneberg-1 extension solver, as the matrix is more dense.

5.1. Synthetic data

For synthetic data, we generate random rank-3 matrices of size $100 \times 100$ with entries uniformly drawn from $[-1, 1]$. All entries are then perturbed with noise drawn from a normal distribution $N(0, \sigma)$ where $\sigma$ takes $(0, 10^{-3}, 10^{-2})$. The rank constraint is enforced using a truncated SVD.

A structured data pattern is formed by removing some measurements in the generated matrices. We are especially interested in the band-diagonal structure which appears in many structure-from-motion problems. A sparse matrix $D$ is said to be band-diagonal with bandwidth $k$ if the following equations hold

$$d_{ij} = 0 \quad \text{for} \quad j < i - k \quad \text{or} \quad j > i + k. \quad (17)$$

Here we generate random band-diagonal matrices of size $100 \times 100$ with varied bandwidth $k$ from 20 down to 4. The corresponding proportion of missing data ranges from 64% up to 92%. For comparison, we consider two state-of-the-art methods, namely Truncated Nuclear Norm Regularization (TNNR-ADMM) [13] and OptSpace [17]. OptSpace was initialized using truncated SVD.

We plot the $\log_{10}$ error versus the bandwidth of the matrix in Fig. 6, with the error defined as $\| W \circ (X - U^T V) \|_F$. When the bandwidth decreases (the missing data increases), the performance of both TNNR-ADMM and OptSpace are affected, especially when the bandwidth $k < 10$ which corresponds to around 80% missing data. In the noise-free case, our method achieves significantly lower error. With low level noise, that is $\sigma = 10^{-3}$, our method remains stable with respect to the rate of missing data. With medium level noise, $\sigma = 10^{-2}$, our method will also be affected when $K < 10$, but will still perform better compared to TNNR-ADMM and OptSpace.

5.2. Real data

Affine structure-from-motion We evaluate our method on the well-known dinosaur sequence for affine structure-from-motion. The original dinosaur sequence contains 2D point tracks from the projection of in total 3184 3D points...
in 36 cameras. Each 3D point is only visible in a few consecutive views and missing for the rest of the views due to self-occlusion. We consider rank-4 factorization which does not use any initial estimate of the translation, see e.g. [28].

To reduce processing time, we take a subset of 3D points which are visible in at least 8 views. This forms an observation matrix $X$ of size $72 \times 116$ with 87.8% missing data. We first compare with TNRR and OptSpace, neither of which handle outliers. In this case, we divide the measurement matrix into 4 overlapping blocks with varied size, see Fig. 7a. The partition of the measurement matrix is a trade-off between the complexity and the numerical accuracy. For a too large block, the numerical error within the block accumulates for the Henneberg-1 extension solver. If we divide the matrix into too many sub-blocks, it unnecessarily increases the complexity. We ran our method with 1000 iterations followed by a non-linear least square optimization in the final step.

The results are shown in Fig. 7. Both TNRR and OptSpace failed to recover the 2D tracks. We plot the recovered tracks for our method and TNRR in Fig. 7c and 7d respectively. When the data matrix is sparse, the recovered 2D tracks from TNRR are stretched towards the origin (top right corner) of the image.

In the presence of outliers, both TNRR and OptSpace will fail. We thus compare with two $L_1$-norm based methods, namely the Wiberg-$L_1$ in [10] and Regularized $L_1$ augmented Lagrange multiplier method (Reg$L_1$-ALM) in [31]. We add 10% outliers to the measurement matrix that in the range $[-50, 50]$. The only change in our method is that in each iteration we compute the error in $L_1$-norm as $\|W \odot (X - U^T V)\|_1$ instead of the Frobenius norm. We run 1000 iterations with our minimal solvers within each sub-block. The average $L_1$ error for our method is 1.314 pixels and 1.231 for the Wiberg-$L_1$ and 2.223 for Reg$L_1$-ALM.

**Linear shape basis** For non-rigid structure-from-motion, a linear shape basis model is commonly used to model the
shape of a non-rigid object. It assumes that any non-rigid deformation of an object can be represented as a linear combination of a set of shapes. Normally the size of the shape basis is much smaller than either the number of frames or the tracked points, so the measurement matrix containing the point tracks can be factorized into a coefficient matrix and a shape basis matrix.

Two datasets – Book and Hand – from [20] are used. The image points are tracked using a standard Kanade-Lucas-Tomasi (KLT) tracker [22]. Due to occlusions, the tracker fails after a number of frames for a subset of points, which leads to the missing data pattern for the Book dataset shown in Fig. 8. In our experiments, we use a subset of 42 frames with 60 tracked points from the Book and 38 frames with 203 points from the Hand dataset. Following the setup in [20], we seek a rank-3 and rank-5 factorization on the two datasets respectively.

The block partition of both our method and [20] are illustrated in Fig. 8. Contrary to [20], missing data within blocks are handled in our method, giving a more flexible partition with wider coverage of measurements. We run our method with 1000 iterations for each block. To further show that our method is capable to handle random missing data, we conduct a second experiment by randomly adding an extra 10% of missing data in the measurement matrix and run both methods with the same settings as before.

We summarize the error $\|W \odot (X - U^TV)\|_F$ in Table 1. Our method achieves smaller error in both dataset with or without adding extra missing data. When random missing data is added into the blocks, [20] fails with large errors. In Fig. 9, the recovered tracked points are plotted for both datasets using our method with extra missing data. Without extra missing data, both our method and [20] achieve quite similar visual results. The running time for our method is 12s on Book and 28s on Hand. For [20], it is 2.5s on Book and 4.5s on Hand. However, the running time for our method can be further reduced as it is highly parallelizable.

### 6. Conclusion

In this paper, we have introduced theory to characterize and generate the minimal problems of low rank matrix factorization, inspired by the Henneberg extensions from rigid graph theory. We have shown that for rank-1 and rank-2, the proposed Henneberg extensions can generate all minimal problems. We conjecture that using additional Henneberg extensions, all minimal problems can be generated for any given rank. Several solvers are proposed to solve minimal problems using constructive extensions. With block partition and a random sampling scheme, minimal solvers can be used in a number of real applications when the data matrix is sparse and contains outliers.
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References


